Chapter 17: Database System Architectures

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Chapter 17: Database System Architectures

- Centralized and Client-Server Systems
- Server System Architectures
- Parallel Systems
- Distributed Systems
- Network Types
Centralized Systems

- Τρέχουν σε έναν υπολογιστή και δεν αλληλεπιδρούν με άλλα συστήματα
- Γενικής χρήσης υπολογιστικό σύστημα: 1- αρκετές CPUs, ελεγκτές με κοινό bus που παρέχει πρόσβαση σε κοινή μνήμη
- Συστήματα ενός χρήστη (e.g., personal computer or workstation): desk-top unit, single user, usually has only one CPU and one or two hard disks; the OS may support only one user.
- Συστήματα πολλαπλών χρηστών: more disks, more memory, multiple CPUs, and a multi-user OS. Serve a large number of users who are connected to the system via terminals. Often called server systems.
A Centralized Computer System

Diagram of a centralized computer system with components such as CPU, disk controller, USB controller, memory, mouse, keyboard, printer, and monitor.
Server systems satisfy requests generated at \( m \) client systems, whose general structure is shown below:
Database functionality can be divided into:

- **Back-end**: πρόσβαση σε δομές, υπολογισμός και βελτιστοποίηση ερωτημάτων, έλεγχος ταυτοχρονισμού και αποκατάσταση

- **Front-end**: εργαλεία όπως φόρμες, γεννήτριες αναφορών και GUIs

Η διασύνδεση μεταξύ front-end και back-end μέσω SQL ή εφαρμογής
Advantages of replacing mainframes with networks of workstations or personal computers connected to back-end server machines:

- better functionality for the cost
- flexibility in locating resources and expanding facilities
- better user interfaces
- easier maintenance
Server System Architecture

- Server systems can be broadly categorized into two kinds:
  - **transaction servers** παρέχουν σύνδεση όπου οι πελάτες στέλνουν αιτήματα
  - **data servers** χρησιμοποιούνται σε object-oriented database systems. Παρέχουν διασύνδεση στο σύστημα αρχείων για δημιουργία, ενημέρωση, αναζήτηση δεδομένων.
Transaction Servers

- Also called **query server** systems or **SQL server** systems
  - Clients send requests to the server
  - Transactions are executed at the server
  - Results are shipped back to the client.
- Requests are specified in SQL, and communicated to the server through a *remote procedure call* (RPC) mechanism.
- Transactional RPC allows many RPC calls to form a transaction.
- **Open Database Connectivity** (ODBC) is a C language application program interface standard from Microsoft for connecting to a server, sending SQL requests, and receiving results.
- JDBC standard is similar to ODBC, for Java
Transaction Server Process Structure

- A typical transaction server consists of multiple processes accessing data in shared memory.

- Server processes
  - These receive user queries (transactions), execute them and send results back
  - Processes may be **multithreaded**, allowing a single process to execute several user queries concurrently
  - Typically multiple multithreaded server processes

- Lock manager process
  - More on this later

- Database writer process
  - Output modified buffer blocks to disks continually
Transaction Server Processes (Cont.)

- Log writer process
  - Server processes simply add log records to log record buffer
  - Log writer process outputs log records to stable storage.
- Checkpoint process
  - Performs periodic checkpoints
- Process monitor process
  - Monitors other processes, and takes recovery actions if any of the other processes fail
    - E.g., aborting any transactions being executed by a server process and restarting it
Transaction System Processes (Cont.)

- User process
- Server process
- ODBC
- JDBC
- Buffer pool
- Query plan cache
- Log buffer
- Lock table
- Process monitor process
- Lock manager process
- Log writer process
- Checkpoint process
- Database writer process
- Log disks
- Data disks
Transaction System Processes (Cont.)

- Shared memory contains shared data
  - Buffer pool
  - Lock table
  - Log buffer
  - Cached query plans (reused if same query submitted again)

- All database processes can access shared memory

- To ensure that no two processes are accessing the same data structure at the same time, database systems implement **mutual exclusion** using either
  - Operating system semaphores
  - Atomic instructions such as test-and-set

- To avoid overhead of interprocess communication for lock request/grant, each database process operates directly on the lock table
  - Instead of sending requests to lock manager process

- Lock manager process still used for deadlock detection
Data Servers

- Used in high-speed LANs, in cases where
  - The clients are comparable in processing power to the server
  - The tasks to be executed are compute intensive.
- Data are shipped to clients where processing is performed, and then shipped results back to the server.
- This architecture requires full back-end functionality at the clients.
- Used in many object-oriented database systems

Issues:
- Page-Shipping versus Item-Shipping
- Locking
- Data Caching
- Lock Caching
Data Servers (Cont.)

- **Page-shipping** versus **item-shipping**
  - Smaller unit of shipping $\Rightarrow$ more messages
  - Worth **prefetching** related items along with requested item
  - Page shipping can be thought of as a form of prefetching

- **Locking**
  - Overhead of requesting and getting locks from server is high due to message delays
  - Can grant locks on requested and prefetched items; with page shipping, transaction is granted lock on whole page.
  - Locks on a prefetched item can be P{called back} by the server, and returned by client transaction if the prefetched item has not been used.
  - Locks on the page can be **deescalated** to locks on items in the page when there are lock conflicts. Locks on unused items can then be returned to server.
Data Servers (Cont.)

- **Data Caching**
  - Data can be cached at client even in between transactions
  - But check that data is up-to-date before it is used (cache coherency)
  - Check can be done when requesting lock on data item

- **Lock Caching**
  - Locks can be retained by client system even in between transactions
  - Transactions can acquire cached locks locally, without contacting server
  - Server **calls back** locks from clients when it receives conflicting lock request. Client returns lock once no local transaction is using it.
  - Similar to deescalation, but across transactions.
Parallel Systems

- Parallel database systems consist of multiple processors and multiple disks connected by a fast interconnection network.

- A coarse-grain parallel machine consists of a small number of powerful processors.

- A massively parallel or fine grain parallel machine utilizes thousands of smaller processors.

- Two main performance measures:
  - **throughput** --- the number of tasks that can be completed in a given time interval
  - **response time** --- the amount of time it takes to complete a single task from the time it is submitted
Speed-Up and Scale-Up

- **Speedup**: a fixed-sized problem executing on a small system is given to a system which is $N$-times larger.
  - Measured by:
    $$\text{speedup} = \frac{\text{large system elapsed time}}{\text{small system elapsed time}}$$
  - Speedup is **linear** if equation equals $N$.

- **Scaleup**: increase the size of both the problem and the system
  - $N$-times larger system used to perform $N$-times larger job
  - Measured by:
    $$\text{scaleup} = \frac{\text{small system small problem elapsed time}}{\text{big system big problem elapsed time}}$$
  - Scale up is **linear** if equation equals 1.
Speedup

- linear speedup
- sublinear speedup

resources

speed
Scaleup

\[ \frac{T_S}{T_L} \]

linear scaleup

sublinear scaleup

problem size
Batch and Transaction Scaleup

- **Batch scaleup:**
  - A single large job; typical of most decision support queries and scientific simulation.
  - Use an $N$-times larger computer on $N$-times larger problem.

- **Transaction scaleup:**
  - Numerous small queries submitted by independent users to a shared database; typical transaction processing and timesharing systems.
  - $N$-times as many users submitting requests (hence, $N$-times as many requests) to an $N$-times larger database, on an $N$-times larger computer.
  - Well-suited to parallel execution.
Factors Limiting Speedup and Scaleup

Speedup and scaleup are often sublinear due to:

- **Startup costs**: Cost of starting up multiple processes may dominate computation time, if the degree of parallelism is high.

- **Interference**: Processes accessing shared resources (e.g., system bus, disks, or locks) compete with each other, thus spending time waiting on other processes, rather than performing useful work.

- **Skew**: Increasing the degree of parallelism increases the variance in service times of parallelly executing tasks. Overall execution time determined by **slowest** of parallelly executing tasks.
Interconnection Network Architectures

- **Bus.** System components send data on and receive data from a single communication bus;
  - Does not scale well with increasing parallelism.

- **Mesh.** Components are arranged as nodes in a grid, and each component is connected to all adjacent components
  - Communication links grow with growing number of components, and so scales better.
  - But may require $2\sqrt{n}$ hops to send message to a node (or $\sqrt{n}$ with wraparound connections at edge of grid).

- **Hypercube.** Components are numbered in binary; components are connected to one another if their binary representations differ in exactly one bit.
  - $n$ components are connected to $\log(n)$ other components and can reach each other via at most $\log(n)$ links; reduces communication delays.
Interconnection Architectures

(a) bus

(b) mesh

(c) hypercube
Parallel Database Architectures

- **Shared memory** -- processors share a common memory
- **Shared disk** -- processors share a common disk
- **Shared nothing** -- processors share neither a common memory nor common disk
- **Hierarchical** -- hybrid of the above architectures
Parallel Database Architectures

(a) shared memory

(b) shared disk

(c) shared nothing

(d) hierarchical
Shared Memory

- Processors and disks have access to a common memory, typically via a bus or through an interconnection network.
- Extremely efficient communication between processors — data in shared memory can be accessed by any processor without having to move it using software.
- Downside — architecture is not scalable beyond 32 or 64 processors since the bus or the interconnection network becomes a bottleneck.
- Widely used for lower degrees of parallelism (4 to 8).
Shared Disk

- All processors can directly access all disks via an interconnection network, but the processors have private memories.
  - The memory bus is not a bottleneck
  - Architecture provides a degree of fault-tolerance — if a processor fails, the other processors can take over its tasks since the database is resident on disks that are accessible from all processors.
- Examples: IBM Sysplex and DEC clusters (now part of Compaq) running Rdb (now Oracle Rdb) were early commercial users
- Downside: bottleneck now occurs at interconnection to the disk subsystem.
- Shared-disk systems can scale to a somewhat larger number of processors, but communication between processors is slower.
Shared Nothing

- Node consists of a processor, memory, and one or more disks. Processors at one node communicate with another processor at another node using an interconnection network. A node functions as the server for the data on the disk or disks the node owns.
- Examples: Teradata, Tandem, Oracle-n CUBE
- Data accessed from local disks (and local memory accesses) do not pass through interconnection network, thereby minimizing the interference of resource sharing.
- Shared-nothing multiprocessors can be scaled up to thousands of processors without interference.
- Main drawback: cost of communication and non-local disk access; sending data involves software interaction at both ends.
Hierarchical

- Combines characteristics of shared-memory, shared-disk, and shared-nothing architectures.
- Top level is a shared-nothing architecture – nodes connected by an interconnection network, and do not share disks or memory with each other.
- Each node of the system could be a shared-memory system with a few processors.
- Alternatively, each node could be a shared-disk system, and each of the systems sharing a set of disks could be a shared-memory system.
- Reduce the complexity of programming such systems by distributed virtual-memory architectures
  - Also called non-uniform memory architecture (NUMA)
Distributed Systems

- Data spread over multiple machines (also referred to as sites or nodes).
- Network interconnects the machines
- Data shared by users on multiple machines
Distributed Databases

- Homogeneous distributed databases
  - Same software/schema on all sites, data may be partitioned among sites
  - Goal: provide a view of a single database, hiding details of distribution

- Heterogeneous distributed databases
  - Different software/schema on different sites
  - Goal: integrate existing databases to provide useful functionality

- Differentiate between local and global transactions
  - A local transaction accesses data in the single site at which the transaction was initiated.
  - A global transaction either accesses data in a site different from the one at which the transaction was initiated or accesses data in several different sites.
Trade-offs in Distributed Systems

- Sharing data – users at one site able to access the data residing at some other sites.
- Autonomy – each site is able to retain a degree of control over data stored locally.
- Higher system availability through redundancy — data can be replicated at remote sites, and system can function even if a site fails.
- Disadvantage: added complexity required to ensure proper coordination among sites.
  - Software development cost.
  - Greater potential for bugs.
  - Increased processing overhead.
Implementation Issues for Distributed Databases

- Atomicity needed even for transactions that update data at multiple sites
- The two-phase commit protocol (2PC) is used to ensure atomicity
  - Basic idea: each site executes transaction until just before commit, and the leaves final decision to a coordinator
  - Each site must follow decision of coordinator, even if there is a failure while waiting for coordinators decision
- 2PC is not always appropriate: other transaction models based on persistent messaging, and workflows, are also used
- Distributed concurrency control (and deadlock detection) required
- Data items may be replicated to improve data availability
- Details of above in Chapter 22
Network Types

- **Local-area networks** (LANs) – composed of processors that are distributed over small geographical areas, such as a single building or a few adjacent buildings.

- **Wide-area networks** (WANs) – composed of processors distributed over a large geographical area.
Local-area Network

- Application server
- Workstation
- Workstation
- Workstation
- Gateway
- Printer
- Laptop
- File server
WANs with continuous connection (e.g., the Internet) are needed for implementing distributed database systems.

Groupware applications such as Lotus notes can work on WANs with discontinuous connection:

- Data is replicated.
- Updates are propagated to replicas periodically.
- Copies of data may be updated independently.
- Non-serializable executions can thus result. Resolution is application dependent.